

Projective coordinates leak

Let us assume that:

- 1) The result of scalar multiplication:

$$P = k \cdot G$$

is available to the adversary in projective coordinates. E.g., by a side-channel attack the adversary is able to find

$$(X, Y, Z),$$

which is an internal representation of P in projective coordinates.

- 2) For simplicity assume that $k \cdot G$ is calculated with the

double-and-add algorithm.

However, the attack works for other methods as well.

The double-and-add method:

Let $(k_{l-1} k_{l-2} \dots k_0)_2$ be the binary representation of k , and $k_l = 1$.

The result $P = k \cdot G$ is calculated as follows:

$$P := G$$

for $(i := l-1; i \geq 0; i--) \{$

$$P := 2 \cdot P$$

$$\text{if } (k_i == 1) \quad P := P + G$$

$\}$
return P .

- 3) Let us fix our attention on the Jacobian coordinates. The standard projective coordinates

shall be summarized later.

The curve $E_{a,b}$ is defined over \mathbb{F}_p , where p is prime.

The aim of the attack is to find the last t bits of k .

Each candidate sequence of bits

$$\underbrace{k_{t-1} k_{t-2} \dots k_1 k_0}_{t \text{ bits}}$$

corresponds to a sequence of operations: point doubling and point addition.

Let s_j be a sequence of intermediate points determined by the operations:

$$s_j = \{ M_u^{(j)} \rightarrow M_{u-1}^{(j)} \rightarrow \dots \rightarrow M_0^{(j)} \}$$

where $M_0^{(j)} = P$.

The arrows above denote additions of G to P or doublings.

We will try to reverse the sequence s_j , that is to move from M_0 to M_1 , further to M_2 , etc., until t -bitting bits of k are revealed.

Sometimes we are successful.

For each step:

$$M_{i+1} \rightarrow M_i$$

we have two possibilities at most:

a) $M_{i+1} \rightarrow M_i$ is an addition

b) $M_{i+1} \rightarrow M_i$ is a doubling.

but if $M_{i+1} \rightarrow M_i$ is an addition then (double and add method) $M_{i+2} \rightarrow M_{i+1}$ must be a point doubling.

Ada When $M_{i+1} \rightarrow M_i$ is an addition For some values (*)
 then, for

$$M_i = (X_i, Y_i, Z_i)$$

$$M_{i+1} = (X_{i+1}, Y_{i+1}, Z_{i+1})$$

$$G = (X_G, Y_G, Z_G) = (x_G, y_G, 1)$$

↑↑
 affine
 coordinates
 of G

$$Z_i = Z_G \cdot Z_{i+1} (X_{i+1} Z_G^2 - X_G Z_{i+1}^2) =$$

$$= 1 \cdot Z_{i+1} (X_{i+1} - x_G \cdot Z_{i+1}^2) =$$

$$= Z_{i+1}^3 (x_{i+1} - x_G)$$

↑ affine $x_{i+1} = \frac{X_{i+1}}{Z_{i+1}^2}$

↑ we are able to calculate this
 as a correct affine
 coordinate of $P-G$.

$$\underbrace{Z_i \cdot (x_{i+1} - x_G)^{-1}}_{(*)} = Z_{i+1}^3$$